Douglas-Rachford; a very versatile algorithm

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Outline I

Feasibility Problems and projection methods

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- Relaxations
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Convex constraints - evolution

of a theory

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References

Much of this presentation is based on material from:

 Scott Lindstrom and BS, "Sixty Years of Douglas-Rachford," (2018); in preparation for the special issue of J. Aust. Maths. Soc. dedicated to Jonathan M. Borwein - as is this talk.







Feasibility Problems and projection methods Convex constraints - evolution of a theory Connection with ADMM Nonconvex Setting

2 sets One application in more detail: ODE BVPs References

2-set feasibility problems

- For two closed constraint sets A and B in a Hilbert space H, find a feasible point $x \in A \cap B$
- When the nearest point projection onto each set is readily computed, the application of a projection algorithm is a popular method of solution.
- Alternating projections [AP], introduced by J. von Neumann in 1933, is the oldest such method.
- Another effective method [D-R], on which we focus, was introduced by J. Douglas and H. H. Rachford in 1956.

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The players







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Projection Methods Alternating Projections [AP] vs. Douglas-Rachford [D-R]





Figure: AP; $T_{A,B} := P_B \circ P_A$ Figure: D-R; $T_{A,B} := \frac{1}{2} (Id + R_B R_A)$

• Projection algorithm: from prescribed x_0 iterate $x_{n+1} = T_{A,B}(x_n)$, then (hopefully) $x = P_A(\{weak-\} \lim_n x_n) \in A \cap B$.

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AP versus D-R





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Can Or Can't

- The efficacy of these methods depends on the ease with which projections onto the constraint sets can be computed,
- HOWEVER:

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Cutter methods - using projections to a separating hyperplane



 In cases where projection onto the set S is difficult to compute we may instead try projecting onto a separating hyperplane, H_S.

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Subgradient Projections



In the special case when the constraint set S is a level set of a convex function f we can try projecting onto the supporting hyperplane to epi(f) at (x₀, f(x₀)).

2 sets Cutter methods Relaxations N sets

Reflection Parameters (relaxed projections)

We could also consider using more general relaxed projections: $R_{H_S}(x)^{(\gamma)} = (2 - \gamma)(P_{H_S}(x) - x) + x.$



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Or, relaxing the averaging step; or both - but we won't.

•
$$T_{A,B}^{(\lambda)} := (1 - \lambda)Id + \lambda R_{H_B}R_{H_A}$$

• Note: $T_{A,B}^{(0)} = \text{Id}$
• and $T_{A,B}^{(1)} = R_BR_A$



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From 2 sets to N sets.

• We can use Douglas-Rachford on a feasibility problem involving N sets; $\Omega_1 \dots \Omega_N$, to find $x \in \bigcap_{k=1}^N \Omega_k$ by utilizing Pierra's product space method [21]; that is, by applying the algorithm in \mathcal{H}^N to the two sets

• $A := \Omega_1 \times \cdots \times \Omega_N$

• $B := \{x = (y_1, \dots, y_N) | y_1 = y_2 = \dots = y_N\}$

- Nicknamed "divide and concur" by Simon Gravel and Veit Elser (the latter credits the former for the name) [17].
 - Reflection in A is the "divide" step entailing reflections in each of the individual constraint sets (eminently parallelizable).
 - "Concur" is the step of reflecting in the agreement (diagonal) set *B*.
- Other methods include cyclically anchored variant (CADRA)
 [5] and Borwein-Tam method (cyclic D-R) [10].

Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

Zeros of sums of maximally monotone operators

Theorem 1 (Lions & Mercier)

In [20] Lions and Mercier showed that when A, B are maximal monotone operators on a H, and A + B is also maximal monotone, then for all non-zero λ and

$$T_{\mathcal{A},\mathcal{B}} \times := J_{\mathcal{B}}^{\lambda} (2J_{\mathcal{A}}^{\lambda} - I) \times + (I - J_{\mathcal{B}}^{\lambda}) \times, \qquad (1)$$

the sequence of iterates, $x_{n+1} = T_{\mathcal{A},\mathcal{B}}x_n$, converges weakly to some $v \in H$, such that $J^{\lambda}_{\mathcal{A}}v$ is a zero of $\mathcal{A} + \mathcal{B}$. Here, $J^{\lambda}_{\mathcal{F}} := (\mathrm{Id} + \lambda \mathcal{F})^{-1}$ is the resolvent operator for $\lambda \mathcal{F}$

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P-L Lions and B. Mercier





Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

View from convex optimization

• If $\mathcal{A} := \partial f$ and $\mathcal{B} := \partial g$ for convex functions f, g, then $0 \in (\partial f + \partial g)u$ means u solves the minimization problem

Find $u \in \operatorname*{argmin}_{x \in X} (f + g)(x)$.

• If, additionally, $f := \iota_A$ and $g := \iota_B$ are indicator functions; recall

$$\iota_c: x \mapsto \begin{cases} 0 & \text{if } x \in C \\ \infty & \text{otherwise,} \end{cases}$$

then $u \in \operatorname{argmin}(f + g)$ implies $u \in A \cap B$ (assuming feasibility).

Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

How a minimization algorithm becomes a projection algorithm

- Phrasing the feasibility problem for A and B as a minimality problem in this way leads us to apply Lions-Mercier's result with $\mathcal{A} := \partial f$ and $\mathcal{B} := \partial g$ where $f := \iota_A$ and $g := \iota_B$ in which case,
 - **(**) $\partial f = N_A$, the normal cone operator

$$x \mapsto egin{cases} \{y | (y, a - x) \leq 0 ext{ for all } a \in A\} & ext{if } x \in C \\ \emptyset & ext{otherwise,} \end{cases}$$

$$\begin{array}{l} \textcircled{0}{2} & \partial g = N_B, \\ \textcircled{0}{3} & J_{\mathcal{A}} := (\mathrm{Id} + \mathcal{A})^{-1} = (\mathrm{Id} + \partial_f)^{-1} = (\mathrm{Id} + N_A)^{-1} = P_A, \\ \textcircled{0}{3} & J_{\mathcal{B}} := (\mathrm{Id} + \mathcal{B})^{-1} = (\mathrm{Id} + \partial_g)^{-1} = (\mathrm{Id} + N_B)^{-1} = P_B, \end{array}$$

• and in Lions-Mercier's algorithm,

$$T_{\mathcal{A},\mathcal{B}} = J_{\mathcal{B}}^{\lambda}(2J_{\mathcal{A}}^{\lambda} - I) + (I - J_{\mathcal{A}}^{\lambda}) = \frac{1}{2}(Id + R_{B}R_{A})$$

Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

Later developments

For $T_{A,B} = \frac{1}{2}R_BR_A + \frac{1}{2}\text{Id} = \frac{1}{2}(2P_B - \text{Id}) \circ (2P_A - \text{Id}) + \frac{1}{2}\text{Id}$:

- Bauschke, Combettes, and Luke [3] were able to give a direct proof that the iterates weakly converge, avoiding the need that $\partial f + \partial B$ be maximal monotone in the feasibility setting.
- The requirement A + B be maximally monotone was later relaxed outside the feasibility setting by Svaiter [23].

B, C and L

terated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs







Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

Definition 2 (Nonexpansivity conditions)

Let $D \subset H$ be nonempty and let $T : D \rightarrow H$. Then T is

1 *nonexpansive* if it is Lipschitz continuous with constant 1:

$$(\forall x, y \in D) \quad \|T(x) - T(y)\| \le \|x - y\|;$$

firmly nonexpansive if

$$(\forall x, y \in D) ||T(x) - T(y)||^2 + ||(Id - T)(x) - (Id - T)(y)||^2 \le ||x - y||^2;$$

Key facts:

1 Projections onto a closed convex set in a Hilbert space are firmly nonexpansive.

2 T is nonexpansive if and only if $\frac{1}{2}(Id + T)$ is firmly nonexpansive

Iterated resolvent operators (Proximal point algorithms) Feasibility as minimization Key ingredients for the proofs

Fejér Monotonicity

Definition 3 (Fejér Monotone)

Where $S \subset H$ is nonempty, the sequence x_n is said to be Fejér *monotone* with respect to S if

$$(\forall y \in S) \ (\forall n \in \mathbb{N}) \ \|x_{n+1} - y\| \le \|x_n - y\|.$$

See, for example, [2].

Proposition 4

If D is a nonempty subset of H and $T : D \to D$ is nonexpansive with $\operatorname{Fix} T \neq \emptyset$ then the sequence $x_{n+1} = T(x_n)$ with $x_0 \in D$ is Fejér monotone with respect to $\operatorname{Fix} T$.

Connection of D-R with ADMM through Duality

• Where F and G are convex, proper, lsc and B is linear, consider the primal problem

$$\mathbf{p} := \inf_{v \in V} \left\{ F(Bv) + G(v) \right\}.$$

 Under sufficient qualification conditions, we can solve p by solving

$$\mathbf{d} := \inf_{v^* \in V^*} \left\{ G^*(-B^*v^*) + F^*(v^*) \right\}.$$

See, for example, Borwein's & Lewis' *Convex Analysis and Nonlinear Optimization* [7, thm 3.3.5].

• Applying DR to **d** is equivalent to applying Uzawa's alternating direction method of multipliers [ADMM] to **p**. See Gabay's 1983 book chapter [16] and the survey [19].

History and independent discovery Advantages A sampling of topics to which D-R has been successfully applied

Nonconvex setting: History and independent discovery

- Fienup independently discovered DR for nonconvex feasibility problems (phase retrieval) [15] and it has been popularized by Veit Elser.
- Other names, special instances, and generalizations:
 - Hybrid Input-Output algorithm (HIO), Fienup's variant, the "difference map" [14]
 - Averaged alternating reflections [4]
 - Relaxed reflect-reflect [13]

History and independent discovery Advantages A sampling of topics to which D-R has been successfully applied

Advantages



DR and AP for a doubleton $A = \{a_1, a_2\}$ and line B in \mathbb{R}^2

- Over-relaxed projection methods tend to explore the space more than AP
- Out-of-the-box solver: need only to be able to compute projections and can apply to any constraint satisfaction problem.

History and independent discovery Advantages A sampling of topics to which D-R has been successfully applied

Theory is scarce

- Although often found to work well, theoretical underpinning for projection methods in the presence of nonconvex constraint sets is sadly lacking.
- Projections onto nonconvex sets are often set valued, and need no longer be firmly nonexpansive, or even nonexpansive.
- Local convergence established in certain instances (in particular near isolated feasible points for intersections of curves and hypersurfaces in ℝⁿ) using theory of local asymptotic stability of almost linear discrete dynamical systems, and more globally utilizing Lyapunov functions. See for example; [9], [6], [8], [18], and [11]

History and independent discovery Advantages A sampling of topics to which D-R has been successfully applied

Discrete/Combinatorial settings:

- Latin squares
- sudoku puzzles
- nonograms
- matrix completion
 - Hadamard matrices
 - Rank minimization
 - distance matrices
- matrix decomposition
- Wavelets with constraints
- 3-SAT
- graph coloring
 - edge colorings
 - 8 queens, 3-SAT, Hamiltonian paths
- Bit retrieval
- doubletons and lines (theory)

Connected constraints:

- Phase retrieval
- Intersections of plane curves and roots of functions
- solving nonlinear systems of equations
- Boundary value ODEs
- Regularity and transversality conditions (theory)

History and independent discovery Advantages A sampling of topics to which D-R has been successfully applied

Two applications where the constraints are discrete finite sets

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

Figure: Solving sudoku puzzles -Elser [13]. Image source Wikimedia Commons [?sudokupicture]



Figure: Solving incomplete Euclidean distance matrices for protein reconstruction [1].

BVPs and formulation as feasibility problems Illustrative example

Nonlinear ordinary differential equations with boundary values [18]

• Consider the problem

y'' = f(y', y, t) for $a \le t \le b$ with $y(a) = \alpha, y(b) = \beta$

- For $N \in \mathbb{N}$, let h := (b a)/(N + 1), and $t_i := a + hi$, for $i = 0, 1, \dots, N + 1$.
- So, t_1, \dots, t_N are a set of equally spaced nodes in [a, b].
- Using centered differences we can obtain a numerical approximation to the solution of the BVP by taking $y(t_i) = \omega_i$, where $\omega_0 = \alpha$, $\omega_{N+1} = \beta$ and for $i = 1, \dots, N$

$$rac{\omega_{i-1}-2\omega_i+\omega_{i+1}}{h^2} \;=\; f\left(t_i,\omega_i,rac{\omega_{i+1}-\omega_{i-1}}{2h}
ight)\; ext{eqn(i)}$$

BVPs and formulation as feasibility problems Illustrative example

Reformulation as a feasibility problem

Let

$$\Omega_i := \{ \omega = (\omega_1, \dots, \omega_N) | \omega \text{ satisfies } eqn(i) \}$$

then finding

$$\omega \in \cap_{k=1}^{N} \Omega_k$$

provides an approximate numerical solution to the BVP..

BVPs and formulation as feasibility problems Illustrative example

Computing projections

• For
$$u := (u_0, u_1, \cdots, u_{N+1})$$
 with $u_0 = \alpha$ and $u_{N+1} = \beta$,

• let
$$H_i(u) = -u_{i-1} + 2u_i - u_{i+1} + h^2 f\left(t_i, u_i, \frac{u_{i+1} - u_{i-1}}{2h}\right)$$
,

• then projection $z := P_{\Omega_i}(v)$ satisfies the Lagrangian:

$$H_i(z) = 0, \quad v - z + \lambda \nabla H_i(z) = 0.$$

• so z can be computed by a steepest descent method or by Newton's method.

BVPs and formulation as feasibility problems Illustrative example

Example: Heaviside

We consider

$$y''(x) = \begin{cases} -1 & y(x) < 0\\ 1 & y(x) \ge 0, \end{cases}$$
(3)

with the boundary conditions y(-1) = -1 and y(1) = 1, which admits the unique continuous solution:

$$y(x) = \begin{cases} -\frac{1}{2}x^2 + \frac{1}{2}x & x < 0\\ \frac{1}{2}x^2 + \frac{1}{2}x & x \ge 0. \end{cases}$$
(4)

BVPs and formulation as feasibility problems Illustrative example

Example: Heaviside



Figure: Newton's Method cycles for certain starting points (left) while DR converges (right).

BVPs and formulation as feasibility problems Illustrative example



Figure: Relative error and error from true solution for converging DR iterates for an ellipse and line.

- The similarities to the previous figure are unmistakable,
- illustrating a frequent feature of the behaviour of D-R.

BVPs and formulation as feasibility problems Illustrative example

Thank You!







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